

Homework 9

Due *Friday*, December 6, 2019

Math 206

1. Write a function `gcd(a,b)` to compute the greatest common divisor of two integers a and b , following the *Euclidean algorithm* that we discussed in class:
 - (a) If $a < b$, switch them so that $a \geq b$.
 - (b) Compute `r = a % b`.
 - (c) If $r = 0$ then b divides a , so return b .
 - (d) Otherwise, replace a with b and b with r , and keep going.

Check that your function gives the right answer for various choices of a and b .

2. Fermat's little theorem states that if p is prime and a is an integer not divisible by p , then

$$a^{p-1} \equiv 1 \pmod{p}.$$

We will discuss the proof on Wednesday, but for now:

- (a) Search online for a list of 3- or 4-digit primes, and do a little experiment to satisfy yourself that this is true.
- (b) Satisfy yourself that it is *not* true when p is not prime.
- (c) You may find that when your numbers get big, it's slow to compute `x**y % z`. Rewrite your code using the Python function `pow(x,y,z)`, which does the same thing much faster.

3. This is a warm-up for *RSA encryption*.

- (a) Choose a 3-digit prime number p and a 4-digit prime number q , and let $n = pq$ and $k = (p - 1)(q - 1)$. Choose a small number e such that $\gcd(e, k) = 1$, and write some code to find a number d such that

$$de \equiv 1 \pmod{k}.$$

- (b) Let $m = 1234$, where m stands for message. Find c , which stands for ciphertext, such that $c \equiv m^d \pmod{n}$.
- (c) The original message should satisfy $m \equiv c^e \pmod{n}$. Check that this is true, and if it's not then find and correct your mistake.

4. Choose a partner with whom to exchange secret messages.

- (a) Write a short secret message. Break it into blocks of 3 letters, and encode each one as a 6-digit number using the code

A = 01	B = 02	C = 03	D = 04	E = 05	F = 06	G = 07
H = 08	I = 09	J = 10	K = 11	L = 12	M = 13	N = 14
O = 15	P = 16	Q = 17	R = 18	S = 19	T = 20	U = 21
V = 22	W = 23	X = 24	Y = 25	Z = 26	period = 27	
					space = 00.	

You could write a function to encode and decode messages, or you could do it by hand.

- (b) Encrypt your 6-digit numbers using $c \equiv m^d \pmod{n}$. Send an email to your colleague containing n and e (your “public key”) and the sequence of 6- or 7-digit numbers c . But do not send p , q , k , or d (your “private key”), nor the unencrypted messages m .
- (c) Have your partner decrypt the message using $m \equiv c^e \pmod{n}$, and decode it using the code above.
- (d) Have your partner write a secret message in reply, encode it as sequence of 6-digit numbers, and encrypt it using *your* public key: that is, $c \equiv m^e \pmod{n}$. Have them send it to you, and decrypt it using your private key: $m \equiv c^d \pmod{n}$. Then decode it and see what they said.

(Note that the roles of d and e are switched in this last part. Your partner still doesn't know your private key.)

5. Do it again with a 6-digit prime for p and a 7-digit prime for q . Probably there is a bottleneck around computing d . Next week we'll discuss a faster way to compute d , using the *extended Euclidean algorithm*. If you get to this problem before then, either think about a clever way to find d , or go read about that algorithm.

For real-life encryption, like sending your credit card number to a web site, we use primes that are hundreds of digits long. Can your code handle primes that big?