## Some large trivalent graphs having small diameters

William M. Kantor\*

Mathematics Department, University of Gregon, Eugene, OR 97403, USA

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Abstract

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If  $n \ge 10$ , then there is a trivalent Cayley graph for G = PSL(n, q) whose diameter is  $O(\log |G|)$ .

This paper concerns an improvement of a result of Babai, Kantor and Lubotzky [1]. In that paper it was shown that there is a constant C such that every non-Abelian finite simple group G has a set S of seven generators for which  $d(G, S) \le C \log |G|$ . Here, S was a carefully chosen generating set for G, and d(G, S) denotes the diameter of the corresponding undirected Cayley graph. This bound is best possible, since a simple count (the "Moore bound") shows that  $d(G, S) + 1 > \log_{2|S|}(|G|)$ .

In this paper we will decrease |S| so as to have |S| = 2 and  $|S \cup S^{-1}| = 3$  in case G = PSL(n, q) with  $n \ge 10$ :

**Theorem.** If  $n \ge 10$ , then there is a trivalent (undirected) Cayley graph for G = PSL(n, q) whose diameter is  $O(\log |G|)$ .

Moreover, there is an algorithm which, when given  $g \in G$ , finds a word in S representing g in  $O(\log |G|)$  steps (i.e., multiplications and inversions of elements of S). Actually, we will only need to assume that  $n \ge 8$  when q is even. There are analogous results obtainable by similar arguments for all the finite simple groups of Lie type, provided that the ranks are not too small. Steinberg [2] obtained two generators for each finite group of Lie type; but his generators do not include an involution, and his argument does not produce the desired diameter bound.

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**Proof.** Given a generating set S, the diameter d(G, S) of the corresponding Cayley graph can be interpreted group-theoretically as the maximum of the lengths of the elements of G as words in  $S \cup S^{-1}$ . We will work inside of SL(n, q), where q is a power of a prime p. In order to obtain a trivalent graph we will find a set  $S = \{s, g\}$  consisting of two matrices, one of which has order 2, such that the corresponding diameter is  $O(\log |G|)$ .

For  $1 \le i,j \le n$  with  $i \ne j$  let  $x_{ij}(\alpha)$  be the matrix with 1's on the diagonal, (i,j)-entry  $\alpha \in \mathbb{F}_q$ , and 0's elsewhere. Then  $X_{ij} := \{x_{ij} \mid \alpha \in \mathbb{F}_q\}$  is isomorphic to the additive group of  $\mathbb{F}_q$ ,  $U := \langle X_{ij} \mid 1 \le i < j \le n \rangle$  is the group of all upper triangular matrices with 1's on the diagonal, and  $U = \prod_{i < j} X_{ij}$  with the  $\frac{1}{2}n(n-1)$  factors written in any order. If  $e_1, \ldots, e_n$  is the standard basis of  $\mathbb{F}_p^n$ , for  $1 \le i < n$  let  $r_i$  and s be the matrices of the transformations behaving as follows:

$$r_i: e_i \rightarrow e_{i+1} \rightarrow -e_i$$
 and  $e_j r_i = e_j$  for  $j \neq i$ ,  $i+1$ ,

and

$$s: e_1 \rightarrow e_2 \rightarrow \cdots \rightarrow e_n \rightarrow (-1)^{n+1}e_1$$
.

Then  $r_{i+1} = r_1^{s^i}$  (where  $g^h := h^{-1}gh$  in any group). If  $t \in \mathbb{F}_q^*$  write  $h_1(t) := \operatorname{diag}(t^{-1}, t, 1, \dots, 1), \ h_{i+1}(t) := h_1(t)^{s^i}$  and  $H_i := \langle h_i(t) | t \in \mathbb{F}_q^* \rangle$  for  $1 \le i < n$ , so that  $H := \prod_i H_i$  is the group of all diagonal matrices in  $\operatorname{SL}(n, q)$ . Also let  $d_1 := \operatorname{diag}(-1, 1, \dots, 1)$  and  $d_{i+1} := d_1^{s^i}$ ; note that  $\det d_i = -1$  and  $d_i^2 = 1$ .

Calculating with  $2\times 2$  matrices, we find that (for any  $t\neq 0$  and  $\alpha$ )

$$x_{i,i+1}(\alpha)^{h_i(t)} = x_{i,i+1}(\alpha t^2), \qquad h_i(t)^{r_i} = h_i(t)^{-1}, \qquad r_i^{d_i} = r_i^{-1} \quad \text{and} \quad r_i^4 = 1.$$

Let  $\theta$  denote a generator of  $\mathbb{F}_q^*$ .

Case 1: q is odd and  $n \ge 12$ . Write  $g := r_1 d_1 \cdot h_3(2) r_3 d_3 \cdot h_5(2\theta) r_5 d_5 \cdot d_7 \cdot x_{9,10}(1) d_9 =$ 

$$\begin{pmatrix}
\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} & 0 & 0 & 0 & 0 \\
0 & \begin{pmatrix} 0 & 1/2 \\ 2 & 0 \end{pmatrix} & 0 & 0 & 0 & 0 \\
0 & 0 & \begin{pmatrix} 0 & 1/2\theta \\ 2\theta & 0 \end{pmatrix} & 0 & 0 \\
0 & 0 & 0 & \begin{pmatrix} -1 & 0 \\ 0 & 1 \end{pmatrix} & 0 \\
0 & 0 & 0 & 0 & \begin{pmatrix} -1 & 1 \\ 0 & 1 \end{pmatrix} & I
\end{pmatrix}$$

We will show that  $S := \{s, g\}$  behaves as required.

Clearly, det g=1 and  $g^2=1$ . In particular,  $|S \cup S^{-1}|=3$ .

**Claim 1.** All elements of  $x_{34}(\mathbb{F}_p)$  have length  $O(\log p)$ .

For,

$$g' := gg^{s^2} = r_1 d_1 \cdot h_3(2) \cdot h_5(\theta) \cdot h_7(2\theta) r_7^{-1} \cdot x_{9,10}(1) \cdot x_{11,12}(1) d_{11},$$

$$g'^4 = h_3(16) h_5(\theta^4) x_{9,10}(4),$$

$$[g'^4, g'^{4s}]^{s^{-8}g'} = [x_{9,10}(4), x_{10,11}(4)]^{s^{-8}g'} = x_{9,11}(16)^{s^{-8}g'} = x_{13}(16)^{g'}$$

$$= x_{13}(16)^{r_1 d_1 h_3(2)} = x_{23}(8).$$

Thus,  $x_{34}(8) = x_{23}(8)^s$  has length O(1), while  $x_{34}(8)^{g'} = x_{34}(8 \cdot 2^2)$ . Now, as in [1], use Horner's rule to express an arbitrary  $t \in \mathbb{F}_p$  in the form

$$t = 8(t/8) = 8 \sum_{i=0}^{m} b_i 2^{2i} = (\cdots(8b_m 2^2 + 8b_{m-1})2^2 + \cdots)2^2 + 8b_0$$

where  $m < \log p$  and the  $b_i$  are integers satisfying  $0 \le b_i < 2^2$ . Then

$$x_{34}(t) = (\cdots((x_{34}(8)^{b_m})^{g'}x_{34}(8)^{b_{m-1}})^{g'}\cdots)^{g'}x_{34}(8)^{b_0}$$

has length  $O(\log p)$ , as claimed.

Claim 2. All elements of  $X_{56} = x_{56}(\mathbb{F}_q)$  and  $X_{65}$  have length  $O(\log q)$ .

For, all elements of  $x_{56}(\mathbb{F}_p) = x_{34}(\mathbb{F}_p)^{s^2}$  have length  $O(\log p)$ . If  $a \in \mathbb{F}_p$ , then  $x_{56}(a)^{g'} = x_{56}(a\theta^2)$ . By writing an arbitrary element of  $\mathbb{F}_q$  in the form  $t = \sum_{i=0}^m a_i \theta^{2i}$ , where  $m < \log_p q$  and  $a_i \in \mathbb{F}_p$ , we can proceed as above to see that each element of  $X_{56}$  looks like

$$x_{56}(t) = (\cdots (x_{56}(a_m)^{g'}x_{56}(a_{m-1}))^{g'}\cdots)^{g'}x_{56}(a_0)$$

for some  $t \in \mathbb{F}_q$  and hence has length  $O(\log q)$ . Now conjugate by g in order to obtain the claim.

From this point on the arguments in [1] can be used, essentially verbatim. We will merely outline them; the reader is referred to that paper for the details. First one shows that all elements of  $L_{12} := \langle X_{12}, X_{21} \rangle \cong \operatorname{SL}(2,q)$  have length  $O(\log q)$ , and hence in particular  $r_1$  and all elements of  $H_1$  do. Then so does  $z := sr_1$ . Note that  $U \subset YY^sY^{s^2} \cdots Y^{s^{n-1}}$  where  $Y := X_{12}X_{12}^zX_{12}^{z^2} \cdots X_{12}^{z^{n-2}}$ , and there are cancellations occurring in these products since  $s^k(s^{k+1})^{-1} = s^{-1}$  and  $z^k(z^{k+1})^{-1} = z^{-1}$ . It follows that each element of Y has length  $O(n \cdot \log q)$ , so that each element of U has length  $O(n \cdot n \log q)$ . Each element of  $H = H_1 H_1^s \cdots H_1^{s^{n-2}}$  also has length  $O(n \log q)$ . Moreover, if  $N := \langle H, r_i | 1 \le i < n \rangle$ , then  $H \le N$ , and each element of  $N/H \cong S_n$  has  $\{r_i H | 1 \le i < n\}$ -length O(n) since the involution  $r_i H$  (of S-length  $O(n \log q)$ ) can be identified with the transposition  $(i, i+1) \in S_n$ . Then each element of N has S-length  $O(n^2 \log q) = O(\log |G|)$ , and hence so does each element of G = UNU.

Case 2: q is odd and n = 10 or 11. This time write  $g := h_1(\theta)r_1d_1 \cdot h_3(2\theta)r_3d_3 \cdot d_5 \cdot x_{78}(1)d_7$  and  $S := \{s, g\}$ , and calculate:

$$g' := gg^{s^2} = h_1(\theta)r_1d_1 \cdot h_3(2) \cdot h_5(2\theta)r_5^{-1} \cdot x_{78}(1) \cdot x_{9,10}(1)d_9,$$
  
$$f := [(gg^{s^2})^4]^{s^{-2}} = h_1(16)x_{56}(4),$$

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$$f^{2} = h_{1}(16^{2})x_{56}(8),$$

$$v := f^{s^{4}} = h_{5}(16)x_{9,10}(4),$$

$$f^{-1}f^{v} = x_{56}(4 \cdot 16^{2} - 4).$$

Thus,  $x_{56}(b)$  has length O(1) for some  $b \in \mathbb{F}_p^*$  (i.e.,  $b = 4 \cdot 16^2 - 4$  or 8), and hence so does  $x_{34}(b) = x_{56}(b)^{s^{-2}}$ . Since  $x_{34}(b)^{g'} = x_{34}(4b)$ , as before it follows that all elements of  $x_{34}(\mathbb{F}_p)$  have length O(log p). Then the same is true of  $x_{i,i+1}(\mathbb{F}_p)$  for each i, and hence also of  $[\cdots[[x_{23}(\mathbb{F}_p), x_{34}(1)], x_{45}(1)], \dots, x_{n1}(1)] = x_{21}(\mathbb{F}_p)$  (since n is bounded!). Now  $r_1 = x_{12}(1)x_{21}(-1)x_{12}(1)$  has length O(log p), and then so does  $g'' := gr_1$ , where  $x_{12}(a)^{g''} = x_{12}(a\theta^2)$ . Now proceed as before.

Case 3: q is even. This time let  $g := r_1 \cdot h_4(\theta) r_4 \cdot x_{78}(1)$  and  $S := \{s, g\}$ . Then

$$g' := gg^{s} = r_{1}r_{2} \cdot h_{4}(\theta)r_{4}h_{5}(\theta)r_{5} \cdot x_{78}(1)x_{89}(1),$$

$$(g^{6})^{s^{-6}g} = [x_{78}(1), x_{89}(1)]^{s^{-6}g} = x_{79}(1)^{s^{-6}g} = x_{13}(1)^{g} = x_{23}(1).$$

Thus,  $x_{78}(1) = x_{23}(1)^{s^5}$  and  $gx_{78}(1) = r_1 \cdot h_4(\theta)r_4$  have length O(1), and hence so does  $u := gx_{78}(1)(gx_{78}(1))^{s^3} = r_1 \cdot h_4(\theta) \cdot h_7(\theta)r_7$ . Since  $x_{45}(a)^u = x_{45}(a\theta^2)$  for all a, by using Horner's rule we find that all elements of  $X_{45}$  have length O(log q), and hence so do all elements of  $X_{54} = (X_{45})^g$ . Now proceed as before.  $\square$ 

It should be noted that a major difference between the cases of odd and even q is that, in the former, in order to use the Hornet's rule argument from [1] we needed to have available  $h_i(2)$  in addition to  $h_j(\theta)$  for some i and j. Those elements were introduced by having the additional dimensions.

A very crude estimate for the diameter obtained in the above argument is  $d(G, S) < 10^7 \log |G|$ .

**Remark.** The analogue of the Theorem holds for the groups  $G = A_n$  and  $S_n$ . We will only indicate this here with an example. It is straightforward to use the methods in [1] to modify this in order to handle the general case.

Let  $G = S_n$  with  $n = 2^{k+1} - 1$  and k odd. Identify the set  $X = \{0, 1, ..., 2^k - 2\}$  with  $\mathbb{Z}_{2^k - 1}$ , and let  $X' = \{x' \mid x \in X\}$  be another copy of X. Consider the n-set  $\{\infty\} \cup X \cup X'$  and (letting x range over X) the permutations

$$t: x \leftrightarrow x', \quad \infty \to \infty,$$
  
 $g:=(\infty,0)(x \to ax) \ (x' \to [ax+a-1]'),$ 

where  $a := 2^{(1/2)(k+1)}$  so that  $a^2 \equiv 2 \pmod{2^k - 1}$ . (Note that  $x \to ax$  fixes 0.) We claim that  $S := \{t, g\}$  behaves as required:  $|S \cup S^{-1}| = 3$  and  $d(G, S) = O(\log |G|)$ . First note that

$$g^2 = (x \to 2x) (x' \to [2x+1]')$$

and

$$(g^2)^t = (x \to 2x + 1) (x' \to [2x]').$$

Every  $x \in X$  is the image of 0 by a word w(x) in  $\{g^2, (g^2)^t\}$  of length  $O(k) = O(\log n)$ : using Horner's rule we can write  $x = \sum_{i=0}^{k} a_i 2^i = 0^{w(x)}$  where  $w(x) := (g^2)^{t^{a_k}} (g^2)^{t^{a_{k-1}}} \cdots (g^2)^{t^{a_k}}$  with all  $a_i \in \{0, 1\}$  (cf. [1]). Also,  $g^k = (\infty, 0)$  since k is odd, so that  $(\infty, 0)$  has length  $O(\log n)$ . If  $x \in X$ , then the transposition  $(\infty, x) = (\infty, 0)^{w(x)}$  also has length  $O(\log n)$ . Then the same is true of every transposition  $(\infty, x') = (\infty, x)^t$ ,  $x \in X$ . Since each element of  $S_n$  is a word of length O(n) in the transpositions just constructed, this proves the claim. This time crude estimates yield that  $d(G, S) < 25n \log n$ .

## References

- [1] L. Babai, W.M. Kantor and A. Lubotzky, Small diameter Cayley graphs for finite simple groups, European J. Combin. 10 (1989) 507-522.
- [2] R. Steinberg, Generators for simple groups, Canad. J. Math. 14 (1962) 277-283.