

All About W

①

Properties: 1) W generated by $\{s_\beta \mid \beta \in \Delta\} = S$ simple reflectors.

2) Def: An expression for $w \in W$ is $w = s_1 s_2 \dots s_d \quad s_i \in S$ (they exist)

Minimal length expression = reduced expression = rex. $l(w) = d$ for rex.

Def: $n(w) = \#$ pos roots sent to neg roots = $\# w\Phi^+ \cap \Phi^- = \# \Phi^+ \cap w^{-1}\Phi^-$ these are the pos roots sent to neg roots.

Prop: $l(w) = n(w)$. Ex: $n(1) = 0 = l(1)$ Ex: $n(s_\beta) = 1 = l(s_\beta)$
 $\# \{ \beta \}$

\hookrightarrow moreover, $\Phi^+ \cap w^{-1}\Phi^- = \left\{ \beta_d, s_d(\beta_{d-1}), s_d s_{d-1}(\beta_{d-2}), \dots \right\} \quad \# = d$
if $w = s_1 s_2 \dots s_d$ rex \lll
 $s_1 \downarrow \Phi^-$ $s_1 s_2 \downarrow \Phi^-$ $s_1 s_2 s_3 \downarrow \Phi^-$

A rex tells you which pos roots sent to neg, and in which order. **Once they go neg, they don't come back.**

$\{ \text{rexs for } w \} \leftrightarrow \{ \text{certain "admissible" orders on } \Phi^+ \cap w^{-1}\Phi^- \}$

Ex: sts vs tst Ex: ss, stst non red

3) W acts freely on $\{ \text{bases} \}$. Pf: If $w \neq 1$ then $w = s_1 s_2 \dots s_d$ and $\beta \mapsto \Phi^- \Rightarrow w\Delta \neq \Delta$.

$\implies W \leftrightarrow \{ \text{bases} \} \Rightarrow W$ finite.

4) $\exists!$ longest word w_0 , $w_0\Delta = -\Delta$, $l(w) = n(w) = \# \Phi^+$.

$w_0\Phi^+ = \Phi^-$

Note that w_0 need not be $-\mathbb{I} \in GL(E)$. If $-\mathbb{I} \in W$ then $w_0 = -\mathbb{I}$ s.t. $-\mathbb{I} \cdot \Delta = -\Delta$.
If $\# \Phi^+$ even/odd then $\det w_0 = \pm 1$.

If $\dim E$ even/odd then $\det(-\mathbb{I}) = \pm 1$. Needn't match up. Even when it does...

Ex: sts Ex: A_n $w_0 = \text{~~XXXX~~}$ $E_1 - E_2 \mapsto E_n - E_{n-1} = -(E_{n-1}, E_n)$

Ex: B_n $w_0 = -\mathbb{I}$.

Fact: ~~XXXXXXXXXX~~ $-w_0 = \tau$ where τ is "Dynkin diagram automorphism", on ~~vertices~~ permutation
of Δ preserving $\langle \alpha, \beta \rangle$.
 $E_1 - E_2 \leftrightarrow E_{n-1} - E_n$
 $E_2 - E_3 \leftrightarrow E_{n-2} - E_{n-1}$
 \vdots

When $\tau = \text{triv}$, $\tau = \mathbb{I}$, $w_0 = -\mathbb{I}$.

5) W has a presentation $W = \langle s_1, s_2 \mid s_i^2 = 1, \underbrace{stst}_{m_{st}} = \underbrace{tsts}_{m_{st}} \Leftrightarrow (st)^{m_{st}} = 1 \rangle$ (2)

where $M_{s_\alpha s_\beta} = \begin{cases} 2 & \theta = 90^\circ & A_1 \times A_1 \\ 3 & \theta = 120^\circ & A_2 \\ 4 & \theta = 135^\circ & B_2 \\ 6 & \theta = 150^\circ & G_2 \end{cases}$ braud rels
 $st = ts$ by ∞ $\theta' = 90, 120, 135, 150 = 180 - \theta = \text{angle between } H_\alpha, H_\beta.$

i.e. W is a Coxeter gp.

Rmk: W is crystallographic, i.e. W preserves a lattice e.g. $\Lambda_{rt} = \mathbb{Z} \cdot \Phi.$

$\{\text{Finite Cryst. Cox Gps}\} = \{\text{Weyl gps of root systems}\}$

only very few other finite Cox gps \Rightarrow dihedral $\langle st \mid m_{st} = \text{ord}(st) \rangle, H_3, H_4$
 have some $m_{st} = 5$
 $4 \cos^2(\frac{180^\circ}{5})$ not in $\mathbb{Z}.$

Ex: $S_{n+1} = \langle s_1, s_2, \dots, s_n \mid s_i^2 = 1, s_i s_{i+1} s_i = s_{i+1} s_i s_{i+1}, s_i s_j = s_j s_i \text{ else} \rangle$ $\mathcal{I} = \parallel$ $\mathcal{X} = \times$ $\mathcal{X} \parallel \mathcal{X} = \times \parallel \mathcal{X}$

$s_2 = \parallel \parallel \times \parallel \parallel$
4 crossings

$l(w) = \# \text{ crossings in a red exp}$

$n(w) = \# \text{ inversions} = \# \{(i < j) \mid w(i) > w(j)\}$

Ex: \times $\begin{matrix} \text{Inv} \\ 23 \\ 13 \\ 24 \\ 14 \end{matrix}$ some a $\Phi^+ \cap w^{-1} \Phi^-$ $\varepsilon_2 - \varepsilon_3 \rightarrow \varepsilon_4 - \varepsilon_1$ | a picture is reduced iff NO TWO STRANDS CROSS TWICE
 inversion removed so $n \neq l.$

Exercise: Think about type B.

6) Any two reds for w are related by braud rels (Matsumoto's Thm)
 (a priori, could have to make longer than shorter.)

7) Suppose L semisimple w/ root system (Φ, h^\star) , and V a fid. L -repr
 which is h -diag'le. $\sum_{\alpha \in \Phi} \text{mult}(\alpha) \text{ wts}(V) \subset h^\star.$ Then

a) $\forall \lambda \in \text{wts}(V), \langle \lambda, \alpha \rangle \in \mathbb{Z}.$ \circ If λ maximal in " α -string thru λ " then $\langle \lambda, \alpha \rangle \in \mathbb{Z}_{\geq 0}.$

b) $W \subset h^\star$, preserves the multiset $\text{wts}(V)$

Pf: L -repr gives s_α repr for each $\alpha \in \Phi.$ So s_α preserves $\text{wts}(V).$ \square

Jumping ahead, but need this lemma soon.

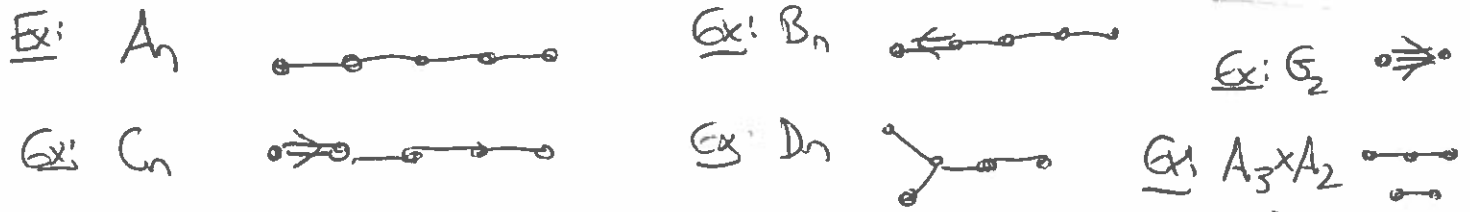
Dynkin Diagrams | Def: Given (Φ, E) and Δ , the Cartan matrix is $(\langle \beta_i, \beta_j \rangle)_{\beta_i, \beta_j \in \Delta}$ (3)

This is NOT the matrix of $(,)$ but, since knowing $\langle \alpha, \beta \rangle$ and $\langle \beta, \alpha \rangle$ determines rank 2 setup, it gives lengths + angles, and determines $(,)$ up to scalar.

Def: The Dynkin diagram is the graph w/ vertices Δ and edges:

$\alpha \rightleftharpoons \beta$ if $\theta = 90^\circ$ (no edge)
 $\alpha \text{---} \beta$ if $\theta = 120^\circ$
 $\alpha \rightrightarrows \beta$ if $\theta = 135^\circ$
 $\alpha \rightleftharpoons\rightleftharpoons \beta$ if $\theta = 150^\circ$

and α is longer.
 Rank: τ gives ord of Γ



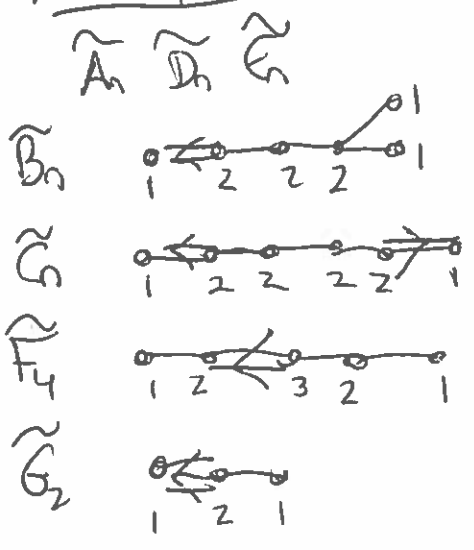
Thm: 1) If (Φ, E) is irred. root system then Γ is either (non-isomorphic list:)




2) Each of these do come from a root system!

PF 1) Similar to first semester. There we already proved: given Γ simply-laced, get pair $(,)$ and ADG only ~~pos def~~ ones. Since pair = Cartan matrix, this agrees. For root simply laced - similar arguments. See Humphreys. Along way, also classify

Affine Dynkin Diagrams: where $(,)$ has a 1D kernel, is pos. semi-def.



To find 1D kernel: label vertices w/ integers d_i s.t.
 $2d_i = \sum \text{neighbors w/ "one-sided mult"}$

\rightleftharpoons is 

PF 2) Construction. See Humphreys. E_8 exercise.
 Note: If you make E_8 then E_6, E_7 are "sub-root systems" i.e. $\text{Span } \{e_i \in \mathbb{Z}\}$.
 Note: All non-simply-laced can be constructed by "folding" simply-laced. Exercise.

Now back to lie algs (eventually back to Lie groups!): We know

$$\{\text{Dynkin Diags}\} \xrightarrow{\cong} \{\text{Irred root systems}\} \xrightarrow{\cong} \{\text{Simple Lie algs}\} \xrightarrow{\cong} \{\text{Simple Lie algs}\}$$

We'll show \Leftrightarrow

(4)

We'll now construct \rightarrow : Given (Φ, E, Δ) well construct L !! By gens + relns

This will make it functional in Dynkin diagrams! Given a sub-root system get a sub-lie-alg!

Given an automorphism τ of Γ get one of L .

Recall: Lie algs by gens + relns. Gens give a v.s. $V = \text{Span}\{\text{gens}\}$.

Get alg $T(V)$. Let $F(V) \subset T(V)$ be subspace spanned by $\{[v_1, [v_2, [v_3, \dots [v_k, v_{k+1}] \dots]]\}$ inside $T(V)$.
This is closed under $[\cdot, \cdot]$ (by Jacobi identity all crazy brackets reduce to this) so is a lie alg.

Free lie alg. Relns: Given a subspace $R \subset F(V)$, let $I \subset T(V)$ be (kernel) ideal generated by R . Then $I(R) \cap F(V)$ is lie alg ideal. $F(V)/I(R) \cap F(V)$ is lie alg by gens + relns.

$$T(V)/I(R) \cong U(\mathfrak{g}(V/R)).$$

Get presentation for U too.

Ex: For \mathfrak{sl}_3 we had basis $\{h_1, h_2, x_1, x_2, x_3, y_1, y_2, y_3\}$ w/ all commutators. This gives
 \rightarrow presentation w/ gens \rightarrow and relns \rightarrow This works for any lie alg!

Ex: For \mathfrak{sl}_3 , $x_3 = [x_1, x_2]$ and $y_3 = [y_1, y_2]$ so don't need as generator.

Reln $[h_i, x_j] = \alpha_j(h_i)x_j$ follows from $[h_1, x_1] = \alpha_1(h_1)x_1$ and $[h_1, x_2] = \alpha_2(h_1)x_2$
and $\alpha_3 = \alpha_1 + \alpha_2$.

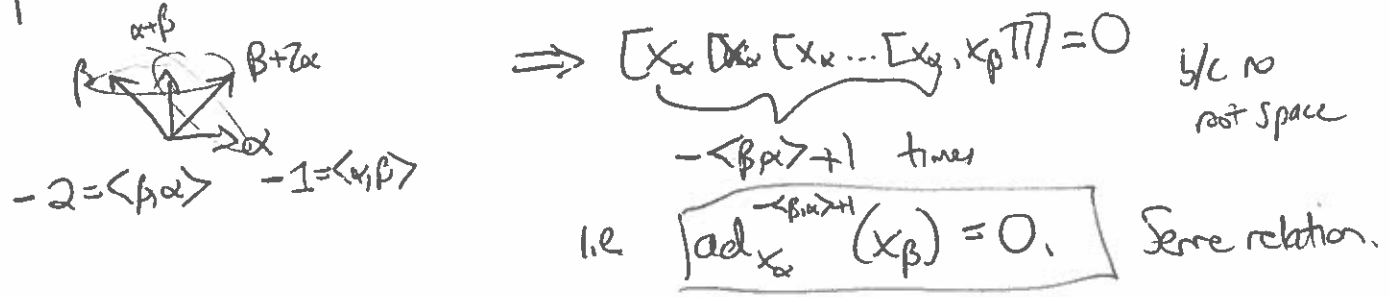
Reln $[x_3, y_3] = h_1 + h_2$ follows from other commutators.

Reln $[x_1, x_3] = 0$ doesn't follow! $[x_1, [x_1, x_2]] = 0$ new relation } Serre Relations.
 $[x_2, x_3] = 0$ $[x_2, [x_2, x_1]] = 0$ \neq

\mathfrak{sl}_3 has presentation: $\langle [x_1, h_1, y_1], [x_2, h_2, y_2] \rangle$ | $[x_i, h_j, y_i]$ \mathfrak{sl}_2 triple, lie diag w/ special espans
 $[x_i, y_j] = 0$ and Serre Relns
 \uparrow $\alpha_i - \alpha_j$ not a root

This presentation is "encoded" in the dynkin diagram!

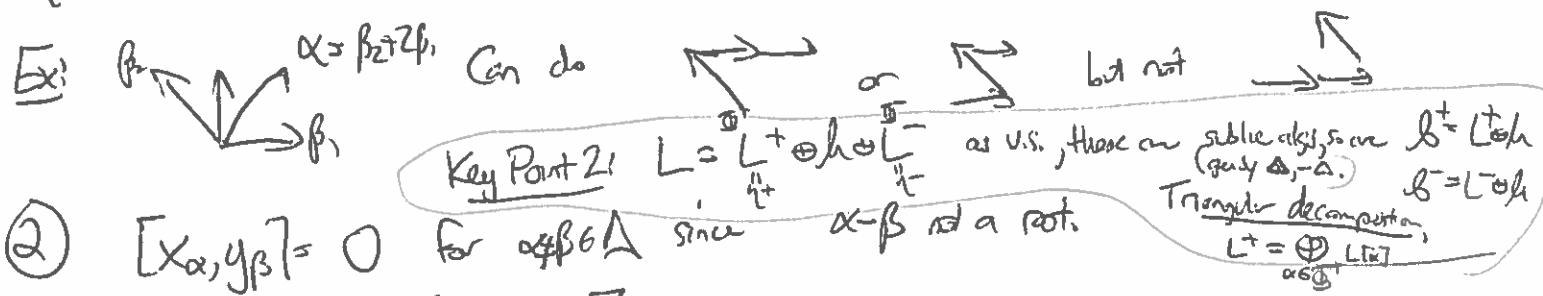
Key point 1: If $\alpha, \beta \in \Delta$ then α -string thru β has length $\langle \beta, \alpha \rangle$ (5)



Minor points: (i) sl_2 triples for $\beta \in \Delta$ do generate all of L . Just x 's generate Φ^+ .

Span $\alpha \in \Phi^+$, $\alpha = \sum c_\beta \beta$ $c_\beta \in \mathbb{Z}_{\geq 0}$. Then can choose a path $\alpha_0 = 0$ $\alpha_1 = \beta_1$ $\alpha_2 = \beta_1 + \beta_2$... $\alpha_n = \alpha$ where each one is a root! (b/c of root string theory)

$\Rightarrow [x_{\beta_1}, [x_{\beta_2}, \dots [x_{\beta_n}, x_\alpha]]] = x_\alpha$ (up to scalar)
 (this is b/c we know $[L[\beta], L[\delta]] = L[\beta + \delta]$ when β, δ are roots.)



(2) $[x_\alpha, y_\beta] = 0$ for $\alpha, \beta \in \Delta$ since $\alpha - \beta$ not a root.

Thm (Serre): A) Fix $(\Phi, \mathfrak{g}, \Delta) \leftrightarrow \Gamma$. Let L_Γ be the Lie algebra defined by gens + rels, where gens are $\{x_\alpha, y_\alpha\}_{\alpha \in \Delta}$ (let $h_\alpha \equiv [x_\alpha, y_\alpha]$) subject to rels: $\mathfrak{h} = \text{span}\{h_\alpha\}$

- (1) $\{x_\alpha, h_\alpha, y_\alpha\}$ sl_2 triple
- (2) $[x_\alpha, y_\beta] = 0$ $\alpha \neq \beta$
- (3) $[h_\alpha, y_\beta] = \langle \beta, \alpha \rangle y_\beta$
 $[h_\alpha, x_\beta] = \langle \beta, \alpha \rangle x_\beta$
(b/c $\beta(h_\alpha)$)
- (4) (S) $(\text{ad } x_\alpha)^{\langle \beta, \alpha \rangle + 1} (x_\beta) = 0$
 $(\text{ad } y_\alpha)^{\langle \beta, \alpha \rangle + 1} (y_\beta) = 0$

Then L_Γ is f.d. and semisimple (simple if Γ connected) w/ max'd toral \mathfrak{h} and root system Φ .
 (b) If L is (semi) simple w/ Φ then $L_\Gamma \cong L$.
Pf: We know $L_\Gamma \twoheadrightarrow L$ since L generated thus at least those rels. So, if L_Γ simple, no ... to claim (A)

Prmk! How do we know $L_P \neq 0$? If L exists w/ P then $L_P \rightarrow L$ is ok. (6)
 Could just construct them explicitly, see Humphreys. But better way is to construct a L_P -repr where it acts faithfully. We'll do that too.

1) Let $P = \text{lie alg by gens + rehs w/o } (S)$. (Expect to be ∞ -dim).

1) $P = P^+ + H + P^-$ P^+ gen by x P^- by y $H = \text{span}\{h\alpha\}$.

This is just b/c of form of rehs. If you take $\vec{x} = [x_1, x_2, \dots, x_{n-1}, x_n]$ and act by $[h, \cdot]$ you rescale ... act by $[y, \cdot]$ you get some h's (which then rescale), etc. Exercise.

(Induction: $[y, \vec{x}] = \underbrace{[y, x_1]}_{\text{Jacobi}} , \vec{x} + \underbrace{[y, \vec{x}]}_{\text{deal with}} , x_1$ unless $n=1$, $[y, x_1] \in H$.
 Induction in P^+

2) H acts semisimply, P graded by Art. $P = P^+ \oplus H \oplus P^-$.

3) More implications: \bullet Any ideal is homogeneous $\bullet P$ is fd. in each degree

- $\bullet \dim P[\alpha] = 1$ for $\alpha \in \Delta$ (just x_α , no more complicated stuff)
- $\bullet \dim P[k\alpha] = 0$ if $k \neq \pm 1$. (b/c $[k\alpha, x_\alpha] = 0$)

\bullet No proper ideal of P intersects H . Pf: If have $\alpha h = h$, have all $\alpha \in P[\Delta]$ with $\chi(h) \neq 0$ (some exist since Δ spans) \Rightarrow have $x_\alpha, \alpha \in \Delta \Rightarrow$ have $h_\beta \Rightarrow$ since P connected, eventually get all generators.

$\bullet P$ has a \downarrow max ll ^{proper} ideal. Pf: Any proper doesn't meet H . Take sum of all proper. Still doesn't meet H . $\Rightarrow P$ has \downarrow simple quotient (if P_H , nonzero!)

4) Constructing an action of P : "Verma module"

Let $\mathbb{C} = \mathbb{C} \cdot v$ be trivial rep of $P^+ \oplus H$, and induce to P to get V .

i.e. $V = U(P^-) \cdot v$ as v.s. $(U(P) = U(P^+) \otimes U(H) \otimes U(P^-)$ as v.s.)

i.e. V has basis $\{y_1, y_2, \dots, y_k v\}$. Action of h is $h \cdot v = 0$

$\Rightarrow h y_1, \dots, y_k v = -\alpha_1(h) y_1, \dots, y_k v + \alpha_2(h) y_1, \dots, y_k v - \dots - \alpha_k(h) y_1, \dots, y_k v = (\quad) y_1, \dots, y_k v.$

$V = \bigoplus V[\lambda]. \quad v \in V[\lambda] \quad y_k$ lowers λ by α_k .

$$x_\alpha \cdot v = 0 \quad \text{so} \quad x_\alpha (y_1 - y_2 v) = \begin{cases} y_1 x_\alpha y_2 - y_2 v & \text{if } y_1 \neq y_2 \\ h_\alpha y_2 - y_2 v & \text{if } y_1 = y_2 \end{cases} \quad (7)$$

both known by induction

These actions satisfy ① ② ③ so P does act on V nontrivially $\Rightarrow P \neq 0$.

Lemma: Let $D \subset P^-$ be ideal generated by $\sum (\text{ad } y_\alpha)^{\langle \beta, \alpha \rangle + 1} (y_\beta)$. Then D is an ideal in P too. (and to paper)

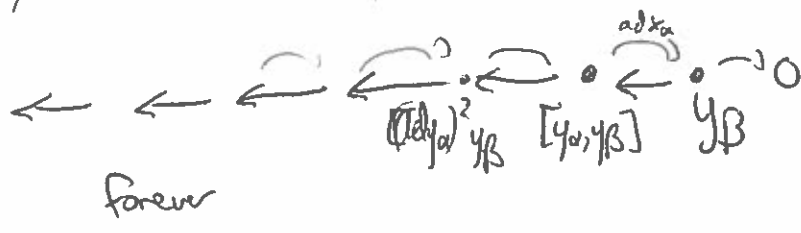
Pf: (This doesn't really use V but philosophically it does) Clearly $[h, D^-] \subset D^-$ since homogeneous. If $\gamma \in \Delta$, $\gamma \neq \alpha, \beta$ then $[x_\gamma y_\alpha] = 0 = [x_\gamma y_\beta]$ so $[x_\gamma S] = 0$.

Now $[x_\beta, y_\alpha] = 0$ so $\text{ad } x_\beta, \text{ad } y_\alpha$ commute, so $[x_\beta, S] = (\text{ad } y_\alpha)^{\langle \beta, \alpha \rangle + 1} ([x_\beta, y_\beta])$

If $\langle \beta, \alpha \rangle \neq 0$ then $[y_\alpha [y_\alpha h_\beta]] = 0$
with y_α

If $\langle \beta, \alpha \rangle = 0$ then $[y_\alpha h_\beta] = 0 \quad \beta \perp \alpha$. ✓

Finally, $[x_\alpha, S] = ?$ Look at α -string of y_β in P (same as in V)



This is ∞ -dim \mathfrak{sl}_2 -rep, given by hw vector of wt $-\beta(h_\alpha) = -\langle \beta, \alpha \rangle \geq 0$.
 we've talked about how $\sum \geq 0$.

we've done this computation before, and exactly

Key Part 3: Same rules are \mathfrak{sl}_2 theory for α -string thru β . $(\text{ad } y_\alpha)^{\langle \beta, \alpha \rangle + 1} (y_\beta)$ is another hw vector.

So D^-, D^+ are \mathfrak{sl}_2 ideals, $D^- + D^+ = (S)$ is an ideal. $L_r \neq 0!!$ $x_\alpha \neq 0, h_\alpha$ lin indep, etc. ETS that L_r is fd simple.

Key Part 4: (S) relations are enough to make $\text{ad } x_\alpha, \text{ad } y_\alpha$ nilpotent on generators !!
 \Rightarrow locally nilp on all of L_r . Iner L_r splits into fd \mathfrak{sl}_2 reps for each triple.

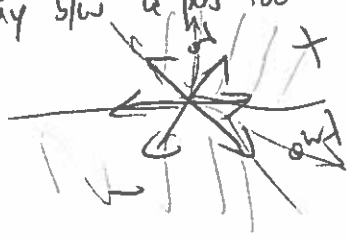
\Rightarrow wts $(L_r) \hookrightarrow W$!!
 Key part 5

But we know $\dim P[\alpha] = \dim L_P[\alpha] = 1 \quad \forall \alpha \in \Delta \Rightarrow \text{true } \forall \alpha \in \Phi$ by Wackin. (8)

$\dim P[k\alpha] = \dim L_P[k\alpha] = 0 \quad \forall \alpha \in \Delta \Rightarrow \text{f.}$

$\dim P[\lambda] = 0$ unless $\lambda = 0$ or $\lambda \in$ positive cone or negative cone (i.e. P^+ or P^-)

but if $\lambda \neq k\alpha$ for some $\alpha \in \Phi$ then λ is on a ray b/w roots, W brings to a ray b/w a pos root + a neg root (not head) so λ not in either cone and $\dim P[\lambda] = 0. \Rightarrow \dim P[\lambda] = 0.$



So $\dim L_P[\alpha] = \begin{cases} 1 & \text{if } \alpha \in \Phi \\ 0 & \text{else} \end{cases}$

$\Rightarrow \mathfrak{h}$ is maxl toral, L_P is simple. \square
w/ root system Φ .

Rank: In similar style, Chevalley constructed algebraic groups (groups in alg geom) for dynkin diagrams. \Rightarrow complex lie grs.
(+ others)
Beyond the scope of this class.